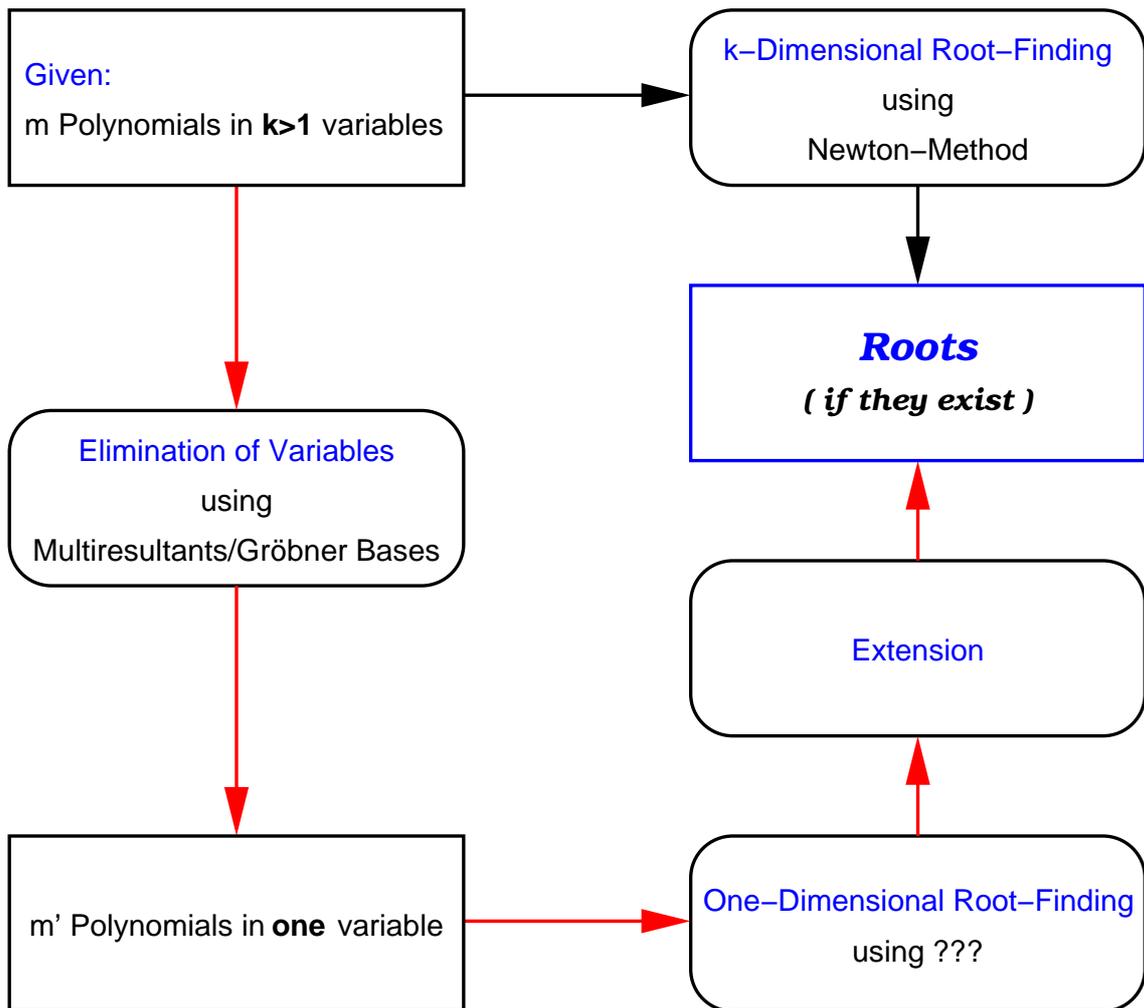


# **Finding Roots of Polynomials**

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# Roots of Polynomials



# Roots of Polynomials in One Variable

## Definition

A **polynomial**  $P$  in one variable  $x$  with complex coefficients is a function, given by

$$P(x) = \sum_{i=0}^n a_i x^i,$$

where  $a_0, \dots, a_n$  are complex numbers with  $a_n \neq 0$ .

## Observation

In general a the roots can be real or complex, single or multiple...

## Theorem 1

If **all coefficients** of  $P$  are **real numbers**, then the complex roots occur in pairs that are conjugate and both roots have the same multiplicity.

# Finding Roots by Polynomial Deflation

## Theorem 2

Every polynomial  $P$  of degree  $n$  with complex coefficients and  $a_n \neq 0$  has the following **product representation**:

$$P(x) = a_n(x - \xi_1)(x - \xi_2) \cdots (x - \xi_n),$$

where  $\xi_1, \dots, \xi_n$  are the roots of  $P$ .

## Idea (Successive **Deflation** of $P$ )

Given a root  $\xi_i$  of  $P$ ,  $1 \leq i \leq n$ , the polynomial can be factored into the following product:

$$P(x) = (x - \xi_i)Q(x).$$

Then the following properties hold:

1. The reduced polynomial  $Q$  has degree one less than  $P$ .
2. The roots of  $Q$  are exactly the remaining roots of  $P$ .

## Remarks

- Deflation is simply **polynomial division**.
- The effort of finding a root hopefully decreases in each step.
- The method cannot converge twice to the same non multiple root.
- Roots become more and more inaccurate, when not **polished up**.
- Successive Deflation is numerical stable, if the root of smallest absolute value is divided out in each step.
- In our context, we don't need complex arithmetics:

$$[x - (a + bi)] \cdot [x - (a - bi)] = x^2 - 2ax + a^2 + b^2 \in \mathbb{R}.$$

# Bracketing

## Definition

A root is **bracketed** in the interval  $(a, b)$  if  $f(a)$  and  $f(b)$  have different signs.

## Motivation

According to the **Intermediate Value Theorem** there must be at least one root in  $(a, b)$ , unless a singularity is present.

## Remark

With **standard arithmetic** there is no sure way of bracketing all roots of an arbitrary function:

$$f(x) = 3x^2 + \frac{1}{\pi^4} \ln[(\pi - x)^2] + 1$$

dips below zero only in the interval  $\pi \pm 10^{-667}$ .

# Bisection

## Idea

- Precondition: A **bracketed range** is given as starting point.
- Evaluate the function at the **midpoint** of the interval and examine its sign.
- Use the midpoint to replace whichever limit has the same sign.

## Remarks

- The method converges **linearly** because the width of the bracketed range decreases by a factor of two after each iteration.
- If the interval contains two or more roots, the method will find only one of them.
- The method does not distinguish singularities from roots.

# Regula Falsi and Secant Method

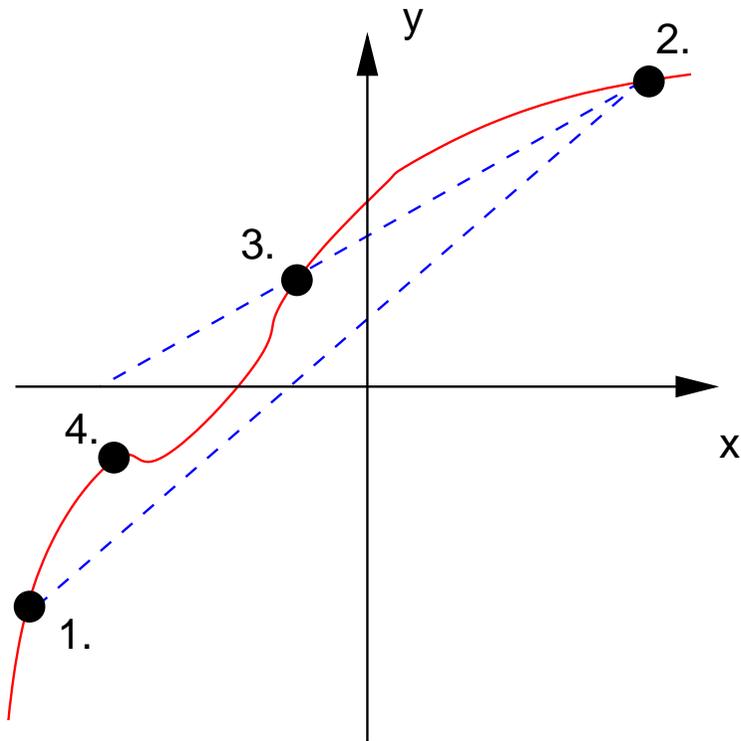
## Assumption

The function is approximately **linear** in the local region of interest.

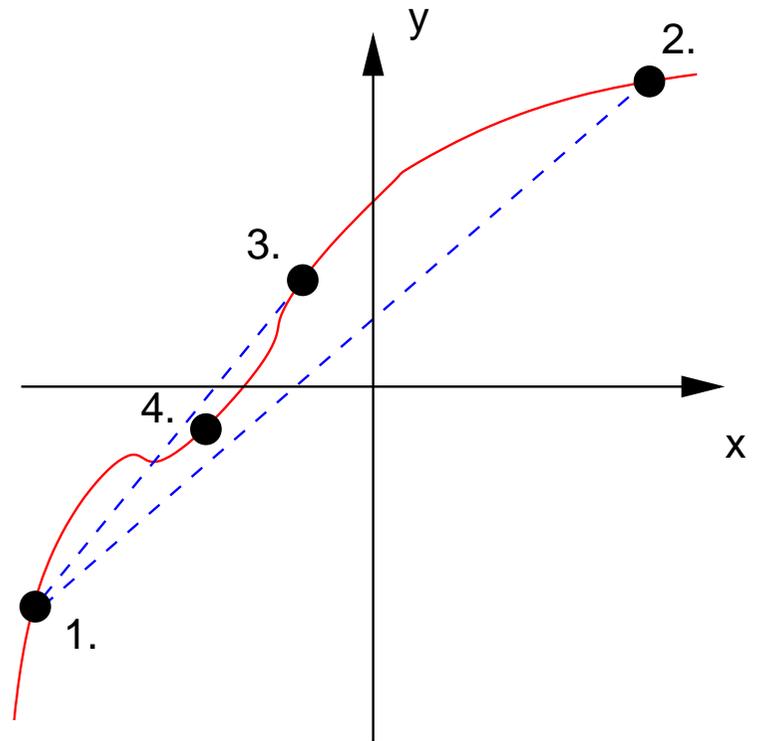
## Idea

- Evaluate the function at the point where the line through both interval limits crosses the axis.
- **Secant Method:** Retain the most recent of the prior estimates and replace the other by the new estimate.
- **Regula Falsi Method:** Retain the prior estimate for which the function value has opposite sign from the current estimate of the root.

# Example



Secant Method



Regula Falsi Method

## Conclusions

- **Secant Method:**
  - The root does not always remain bracketed. There is no convergence guaranty.
  - **Near** the root of a **sufficient continuous** function the convergence order is the "golden ratio"  $1.618\dots$
- **Regula Falsi Method:**
  - Convergence can be guaranteed since the root remains bracketed.
  - Convergence order is lower as in the case of the Secant Method.

## Improvements

- **RIDDER's Method:** Variant of the Regula Falsi Method that uses exponential instead of linear interpolation.
- **BRENT's Method:** Combines Secant Method, Bisection and quadratic interpolation.
- **MULLER's Method:** Generalization of the Secant Method using quadratic interpolation.

## LAGUERRE's Method

### Main Idea

- The root  $\xi_1$  that we seek is assumed to be some distance  $a$  from our current guess  $\hat{\xi}_1$ .
- All other roots are assumed to be located at a distance  $b$ .
- Use the polynomial  $P, P', P''$  to solve for  $a$ , then take  $\hat{\xi}_1 - a$  as the next guess.
- Continue this process until  $a$  becomes small enough.

## Some Details

Remember the following relations between  $P$  and its derivatives:

$$P(x) = \prod_{i=1}^n (x - \xi_i)$$

$$Q(x) := \frac{P'(x)}{P(x)} = \sum_{i=1}^n \frac{1}{x - \xi_i}$$

$$R(x) := \left[ \frac{P'(x)}{P(x)} \right]^2 - \frac{P''(x)}{P(x)} = \sum_{i=1}^n \frac{1}{(x - \xi_i)^2}$$

Using our "rather drastic set of assumptions"

$$\xi_1 = \hat{\xi}_1 - a, \quad \xi_i = \hat{\xi}_1 - b, \quad 2 \leq i \leq n.$$

we obtain for  $Q(\hat{\xi}_1)$  and  $R(\hat{\xi}_1)$ :

$$Q(\hat{\xi}_1) = \frac{1}{a} + \frac{n-1}{b} \quad R(\hat{\xi}_1) = \frac{1}{a^2} + \frac{n-1}{b^2}.$$

Solving for  $a$  leads to:

$$a = \frac{n}{Q(\hat{\xi}_1) \pm \sqrt{(n-1)(nR(\hat{\xi}_1) - Q(\hat{\xi}_1)^2)}}$$

## Remarks

- There are two possibilities for  $a$ . Take the sign, such that  $a$  is minimal.
- For polynomials with **all real roots** the method is guaranteed to converge to a root for any starting point.
- For polynomials with **some complex roots** convergence cannot be guaranteed.
- When the method converges to a simple complex root the convergence is **third order**.
- The method requires complex arithmetic, even while converging to real roots.

# Eigenvalue Methods

## Facts

- The **eigenvalues** of a matrix  $A$  are the roots of the **characteristic polynomial**

$$Q_A(x) = \det(A - xI).$$

- There are efficient and numerical stable **non-root-finding** methods to compute the eigenvalues of a matrix.

## Question

Is it possible to **reduce** the polynomial root-finding problem to the problem of computing the eigenvalues of a matrix ?

## Answer

The characteristic polynomial of the following [companion matrix](#)

$$A = \begin{pmatrix} -\frac{a_{n-1}}{a_n} & -\frac{a_{n-2}}{a_n} & \cdots & -\frac{a_1}{a_n} & -\frac{a_0}{a_n} \\ 1 & 0 & \cdots & 0 & 0 \\ 0 & 1 & \cdots & 0 & 0 \\ \vdots & & \ddots & & \vdots \\ 0 & 0 & \cdots & 1 & 0 \end{pmatrix}$$

has the same roots as the polynomial

$$P(x) = \sum_{i=0}^n a_i x^i.$$

## Proof

Expansion by the first row gives us

$$Q_A(x) = \det(A - xI) = (-1)^n \sum_{i=0}^n \frac{a_i}{a_n} x^i.$$

## Remarks

- The eigenvalues can be computed using the *QR-Algorithm*, an efficient eigenvalue method when the input is an upper Hessenberg matrix.
- *Advantage:*  
More *robust technique* than LAGUERRE'S Method.
- *Disadvantage:*  
Typically a factor 2 *slower* than LAGUERRE'S Method.

## Simultaneous Inclusion of Real Roots

Given:  $P(x) = \sum_{i=0}^n a_i x^i$ ,  $a_n = 1$

### Assumptions

1.  $P$  has  $n$  real roots  $\xi = (\xi^{(1)}, \dots, \xi^{(n)})$ , where multiple roots are entered according to their multiplicity.
2. We collect multiple roots as  $(\xi^{(m+1)}, \dots, \xi^{(n)})$  and forget about them.
3. For all roots we know including intervals

$$X^{(0,i)} = [x_1^{(0,i)}, x_2^{(0,i)}], \quad 1 \leq i \leq m.$$

4. Including intervals are pairwise disjoint:

$$X^{(0,i)} \cap X^{(0,j)} = \emptyset, \quad 1 \leq i < j \leq m.$$

We consider the **equivalent** polynomial  $Q$  where every root has multiplicity 1:

$$Q(x) = \frac{P(x)}{\text{GCD}(P(x), P'(x))} = \prod_{j=1}^m (x - \xi^{(j)}).$$

Extracting  $(x - \xi^{(i)})$  leads to

$$\xi^{(i)} = x - Q(x) / \prod_{j=1, j \neq i}^m (x - \xi^{(j)}).$$

If we choose  $x = x^{(0,i)} \in X^{(0,i)}$ , we have  $\xi^{(i)} \in$

$$X^{(1,i)} := \left\{ x^{(0,i)} - \frac{Q(x^{(0,i)})}{\prod_{j=1, j \neq i}^m (x^{(0,i)} - X^{(0,j)})} \right\} \cap X^{(0,i)}$$

... and the following **total step** iteration scheme:

$$X^{(k+1,i)} := \left\{ x^{(k,i)} - Q(x^{(k,i)}) / B^{(k,i)} \right\} \cap X^{(k,i)}$$

$$B^{(k,i)} := \prod_{j=1, j \neq i}^m (x^{(k,i)} - X^{(k,j)})$$

$$x^{(k,i)} \in X^{(k,i)}, \quad 1 \leq i \leq m, \quad k \geq 0.$$

## Improvements

Before computing  $X^{(k+1,j)}$ ,  $j \geq i$ , we **already** know  $X^{(k+1,j)}$ ,  $j < i$ . Therefore we can replace  $B^{(k,i)}$  by

$$C^{(k,i)} := \prod_{j=1}^{i-1} (x^{(k,i)} - X^{(k+1,j)}) \cdot \prod_{j=i+1}^m (x^{(k,i)} - X^{(k,j)}).$$

to get a tighter denominator and the following **single step** iteration scheme:

**Init:**  $x^{(0,i)} \in X^{(0,i)}$

**Step:**

$$X^{(k+1,i)} := \left\{ x^{(k,i)} - Q(x^{(k,i)})/C^{(k,i)} \right\} \cap X^{(k,i)}$$

$$C^{(k,i)} := \prod_{j=1}^{i-1} (x^{(k,i)} - X^{(k+1,j)}) \cdot \prod_{j=i+1}^m (x^{(k,i)} - X^{(k,j)}),$$

$$x^{(k+1,i)} \in X^{(k+1,i)}, \quad 1 \leq i \leq m, \quad k \geq 0.$$

**Heuristic:** Choose  $x^{(k,i)} = \frac{1}{2}(x_1^{(k,i)} + x_2^{(k,i)})$

## Conclusions

- **Advantages**

- **Simultaneous** determination of polynomial roots.
- **Reliable** information about root location.
- **Always converging** under the assumptions made above.
- **Early sign prediction** of roots possible.
- Convergence order  $\geq 2$  for the total step method and  $> 2$  in the case of the single step method.

- **Disadvantages**

- **Interval arithmetic** is expensive.
- **All** roots of the polynomial have to be **real**.

# Simultaneous Inclusion of Complex Roots

## Assumption

We are using **circular regions** as complex intervals.

The following relationship between  $P$  and  $P'$ :

$$\begin{aligned} Q(z) = \frac{P'(z)}{P(z)} &= \frac{\sum_{i=1}^m \prod_{j=1, j \neq i}^m (z - \xi^{(j)})}{\prod_{j=1}^m (z - \xi^{(j)})} \\ &= \sum_{i=1}^m \frac{1}{z - \xi^{(i)}} \end{aligned}$$

let us write  $\xi^{(i)}$  as follows:

$$\begin{aligned} \xi^{(i)} &= z - (z - \xi^{(i)}) = z - \frac{1}{\frac{1}{z - \xi^{(i)}}} \\ &= z - \frac{1}{Q(z) - \sum_{j=1, j \neq i}^m \frac{1}{z - \xi^{(j)}}}. \end{aligned}$$

This leads to the following **total step** iteration scheme:

$$\begin{aligned}
 Z^{(k+1,i)} &= \langle z^{(k+1,i)}, r^{(k+1,i)} \rangle \\
 &= z^{(k,i)} - \frac{1}{Q(z^{(k,i)}) - C^{(k,i)}} \\
 z^{(k,i)} &= m(Z^{(k,i)}), \\
 Q(z^{(k,i)}) &= \frac{P'(z^{(k,i)})}{P(z^{(k,i)})} \quad \text{for } P(z^{(k,i)}) \neq 0, \\
 C^{(k,i)} &= \sum_{j=1, j \neq i}^m \frac{1}{z^{(k,i)} - Z^{(k,i)}}, \\
 &1 \leq i \leq m, \quad k \geq 0.
 \end{aligned}$$

## Remark

- A **single step** method can be formulated as in the case of real roots.
- The convergence is at least **third order**.